

Chapter 23

Sorting

Selection Sort

- **Strategy** (Ascending, Array)
 - Select the smallest element, place in proper position
 - Repeat step 1 with remaining items
- Does not depend on the initial arrangement of the data
- Only appropriate for small n

Selection Sort :: Code

```
public static void selectionSort(int[] list) {
    for (int top = 0; top <= list.length - 2; top++) {
        swap(list, top, findMinIndex(list, top));
    }
}
```

```
private static int findMinIndex(int[] list, int startAt) {
    int minIndex = startAt;
    for (int index = startAt + 1; index <= list.length - 1; index++) {
        if (list[index] < list[minIndex]) {
            minIndex = index;
        }
    }
    return minIndex;
}
```

```
private static void swap(int[] list, int index1, int index2) {
    int tmp = list[index1];
    list[index1] = list[index2];
    list[index2] = tmp;
}
```

Bubble Sort

- Strategy (Ascending, Array)
 - Compare adjacent elements ($[i]$ and $[i+1]$) and exchange them if they are out of order ($[i] > [i+1]$)
 - Will sink the largest element to the end of the array
 - Repeating this process will eventually sort the array into ascending order

Bubble Sort :: Code

```
public static void bubbleSort(int[] list) {
    boolean sorted = false;
    int pass = 1;

    while (pass <= list.length - 1 && !sorted) {
        sorted = true;
        for (int i = list.length - 1; i >= pass; i--) {
            if (list[i] < list[i - 1]) {
                swap(list, i, i - 1);
                sorted = false;
            }
        }
        pass += 1;
    }
}
```

Insertion Sort

- **Strategy** (Ascending, Array)
 - Place item $[i]$ where $i = 1$ in its proper location with relation to items $[0] \dots [i-1]$
 - Repeat step 1 with rest of items $[2] \dots [n-1]$

Insertion Sort :: Code

```
public static void insertionSort(int[] list) {
    int holdItem;
    int j;

    for (int i = 1; i <= list.length - 1; i++ ) {
        holdItem = list[i];
        for (j = i - 1; j >= 0 && list[j] > holdItem; j--) {
            list[j + 1] = list[j];
        }
        list[j + 1] = holdItem;
    }
}
```

MergeSort

- A recursive sorting algorithm
- Gives the same performance, regardless of the initial order of the array items
- Strategy
 - Divide an array into halves
 - Sort each half
 - Merge the sorted halves into one sorted array

MergeSort :: Code

```
public static void mergeSort(int[] list) {  
    if (list.length > 1) {  
        int n2 = list.length / 2 + 1;  
        int[] leftHalf = new int[n2];  
        System.arraycopy(list, 0, leftHalf, 0, n2);  
        mergeSort(leftHalf);  
  
        int n = list.length - n2;  
        int[] rightHalf = new int[n];  
        System.arraycopy(list, list.length / 2, rightHalf, 0, n);  
        mergeSort(rightHalf);  
  
        merge(leftHalf, rightHalf, list);  
    }  
}
```

MergeSort :: Code

```
public static void merge(int[] left, int[] right, int[] list) {
    int leftIndex = 0;
    int rightIndex = 0;
    int listIndex = 0;

    while (leftIndex < left.length && rightIndex < right.length) {
        if (left[leftIndex] < right[rightIndex])
            list[listIndex++] = left[leftIndex++];
        else
            list[listIndex++] = right[rightIndex++];
    }

    while (leftIndex < left.length)
        list[listIndex++] = left[leftIndex++];

    while (rightIndex < right.length)
        list[listIndex++] = right[rightIndex++];
}
```

QuickSort

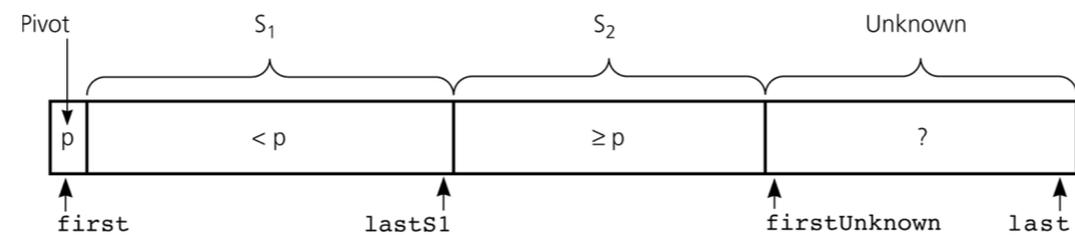
- A divide-and-conquer algorithm
- Strategy
 - Partition an array into items that are less than the pivot and those that are greater than or equal to the pivot
 - Sort the left section
 - Sort the right section

QuickSort: Algorithm

```
QuickSort(array:ArrayType, first:int, last:int) {  
  if (first < last) {  
    Choose a pivot item p from array[first..last]  
    Partition the items of array[first..last] about p  
  
    // the partition is array[first..pivotIndex..last]  
    // sort S1  
    QuickSort(array, first, pivotIndex - 1)  
  
    // sort S2  
    QuickSort(array, pivotIndex + 1, last)  
  }  
}
```

QuickSort: Partition Invariant

- Use the invariant to develop a partition algorithm
 - The items in region S_1 are all less than the pivot, and those in S_2 are all greater than or equal to the pivot



Partition: Algorithm

```
Partition(array:ArrayType, first:int, last:int, pivotIndex:int)
  choose the pivot and swap it with array[first]
  p = array[first]

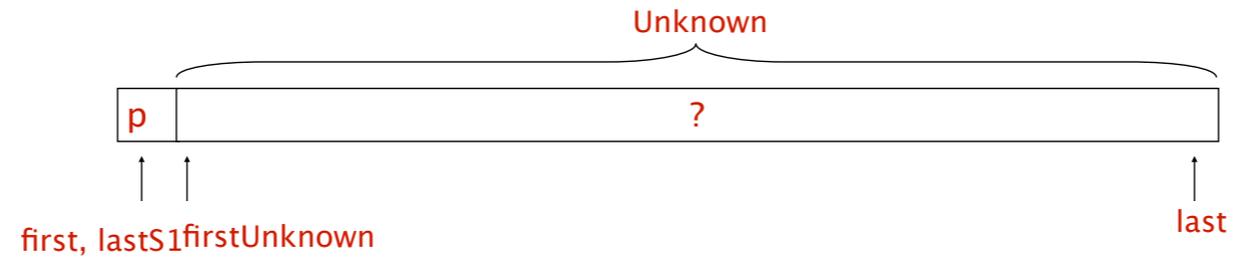
  lastS1 = first
  firstUnknown = first + 1

  while (firstUnknown <= last) {
    if (array[firstUnknown] < p)
      Move array[firstUnknown] into S1
    else
      Move array[firstUnknown] into S2
  }

  swap array[first] with array[lastS1]
  pivotIndex = lastS1
```

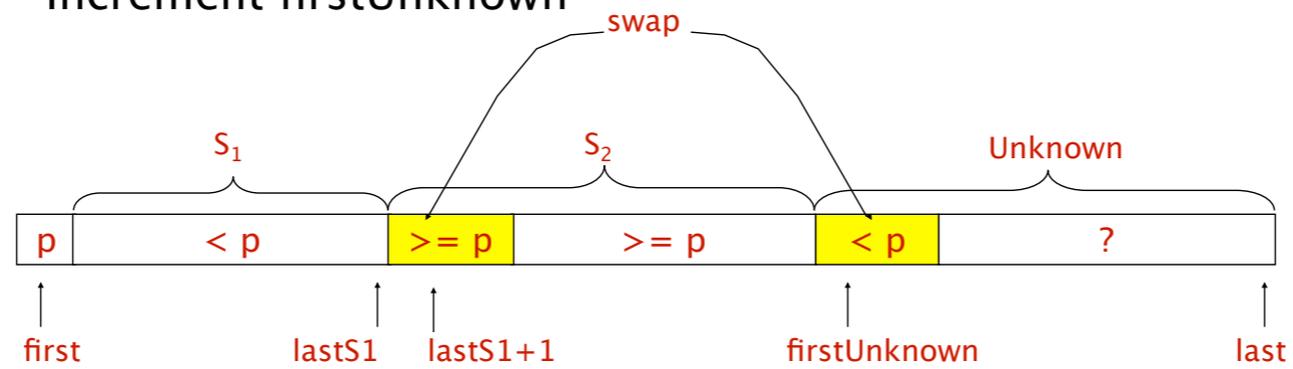
Partition: Initial Array State

```
lastS1 = first  
firstUnknown = first + 1
```



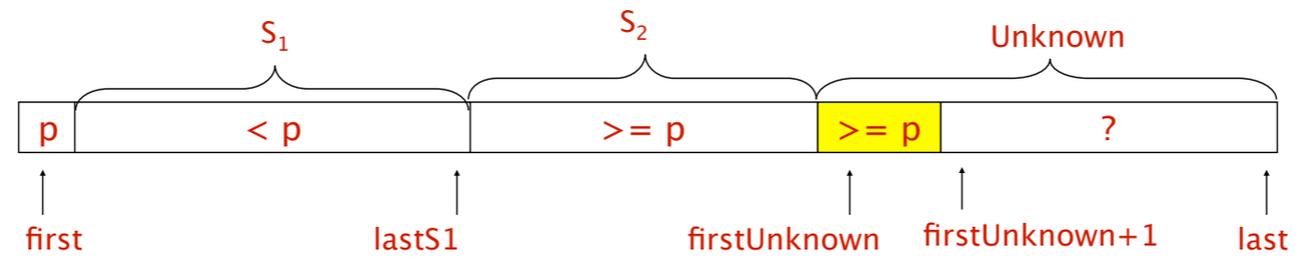
Partition: Move array[firstUnknown] into S_1

swap array[firstUnknown] with array[lastS1 + 1]
Increment lastS1
Increment firstUnknown

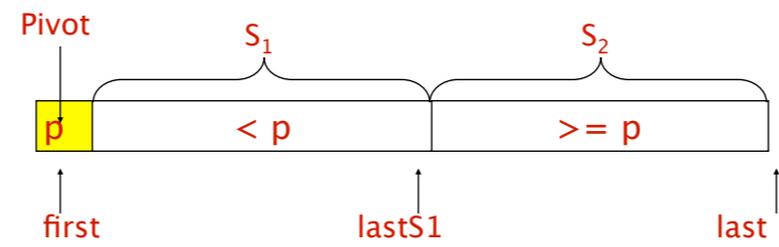


Partition: Move array[firstUnknown] into S_2

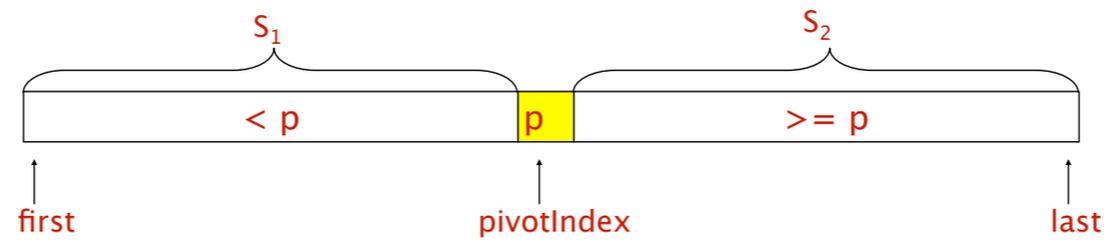
Increment firstUnknown



Partition: Final Task



swap array[first] with array[lastS1]
pivotIndex = lastS1



Heapsort: Strategy

- Transforms the array into a heap
- Removes the heap's root (the largest element) by exchanging it with the heap's last element
- Transforms the resulting semiheap back into a heap

Heapsort

- Compared to mergesort
 - Both heapsort and mergesort are $O(n * \log n)$ in both the worst and average cases
 - However, heapsort does not require a second array
- Compared to quicksort
 - Quicksort is $O(n * \log n)$ in the average case
 - It is generally the preferred sorting method, even though it has poor worst-case efficiency : $O(n^2)$

Radix: Traced

0123, 2154, 0222, 0004, 0283, 1560, 1061, 2150	Original integers
(156 0 , 215 0) (106 1) (022 2) (012 3 , 028 3) (215 4 , 000 4)	Grouped by fourth digit
1560, 2150, 1061, 0222, 0123, 0283, 2154, 0004	Combined
(000 4) (022 2 , 012 3) (215 0 , 215 4) (156 0 , 106 1) (028 3)	Grouped by third digit
0004, 0222, 0123, 2150, 2154, 1560, 1061, 0283	Combined
(000 4 , 106 1) (012 3 , 215 0 , 215 4) (022 2 , 028 3) (156 0)	Grouped by second digit
0004, 1061, 0123, 2150, 2154, 0222, 0283, 1560	Combined
(000 4 , 012 3 , 022 2 , 028 3) (106 1 , 156 0) (215 0 , 215 4)	Grouped by first digit
0004, 0123, 0222, 0283, 1061, 1560, 2150, 2154	Combined (sorted)

Radix Sort: n decimal integers, d digits each

```
RadixSort(array:ArrayType, n:int, d:int)
  for (j = d down to 1) {
    Initialize 10 groups to empty
    Initialize a counter for each group to 0

    for (i = 0 through n-1) {
      k = jth digit of array[i]
      Place array[i] at the end of group k
      Increase kth counter by 1
    }

    Replace the items in array with all items in group 0, followed by all
    the items in group 1 and so on.
  }
```